

# Dynamic Vehicle Routing Problem

## Project Work III.

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# Introduction

- The Dynamic Vehicle Routing Problem (DVRP) involves adapting vehicle routes in real time to handle new delivery requests (e.g., factory logistics, meal deliveries).
- Seeking a solution that is only statically optimal may not always be the best approach, as it may lack the flexibility needed to handle dynamic changes.
- A more adaptive schedule is often required, one that can efficiently accommodate new requests while minimizing delays.

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- Seeking a solution that is only statically optimal may not always be the best approach, as it may lack the flexibility needed to handle dynamic changes.
- A more adaptive schedule is often required, one that can efficiently accommodate new requests while minimizing delays.
- During the semester, I focused on a detailed study of a competition problem, and implement a natural algorithm.
- While effective, the algorithm is resource intensive and lacks explicit adaptability to dynamic changes.

# Huawei Competition

We follow the notation system from [1]. The input consists of the following:

- A directed graph  $G = (F, A)$ , where  $F$  is the set of factories, and  $A$  is the set of arcs connecting the factories. Each arc has a transportation time  $t_{ij}$ .
- An order set  $O = \{o_i : i = 1, \dots, N\}$ , where each order  $o_i = (F_p^i, F_d^i, q^i, t_e^i, t_j^i)$  specifies:
  - $F_p^i$  and  $F_d^i$ : the pickup and delivery locations.
  - $q^i = (q_{\text{standard}}^i, q_{\text{small}}^i, q_{\text{box}}^i)$ : the size of the order in pallets and boxes.
  - $t_e^i$ : the creation time of the order.
  - $t_j^i$ : the committed completion time.
- A fleet of vehicles  $V = \{v_k : k = 1, \dots, K\}$ , each with a loading capacity and specific shift times.
- $M$  nodes (factories), where each factory has limited cargo docks and work shifts. Vehicles may need to wait if all docks are busy.

# Constraints

The problem must satisfy the following constraints:

- 1 **Order fulfillment:** All orders must be served.
- 2 **Completion time:** Orders must be completed before their committed times  $t_j^i$ .
- 3 **Order splitting:** Orders cannot be divided across multiple vehicles unless specified.
- 4 **Vehicle capacity:** No vehicle can exceed its loading capacity.
- 5 **Work shifts:** Loading and unloading must occur within shift times. (*We do not take this restriction into account.*)
- 6 **Dock limitations:** Each factory has limited docks, and vehicles follow a first-come, first-serve rule.

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There are also some hidden constraints in the problem that are not explicitly mentioned in the problem statement, but are assumed during the validation process.

# Objective function

The problem has two main objectives:

- 1 Minimize the total delay of orders (*tardiness*):

$$f_1 = \sum_{i=1}^N \max(0, a_i^d - t_i^l),$$

where  $a_i^d$  is the arrival time of order  $o_i$ ,  $t_i^l$  is the committed completion time, and  $N$  is the total number of orders.

- 2 Minimize the average travel distance of vehicles:

$$f_2 = \frac{1}{K} \sum_{k=1}^K \sum_{i=1}^{l_k-1} d_{n_i^k, n_{i+1}^k},$$

where  $n_i^k$  is the  $i$ -th node in the route of vehicle  $v_k$ ,  $d_{n_i^k, n_{i+1}^k}$  is the distance between consecutive nodes, and  $K$  is the total number of vehicles.

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The overall objective function is:  $f = \lambda \cdot f_1 + f_2$ , where  $\lambda$  is a large positive constant to prioritize minimizing delays. In the validator, it is fixed as

$$\lambda = \frac{10\,000}{3600}$$



# Our work in this semester

- Getting familiar with the area, item Understanding the concepts of my advisors general simulation framework. (I rely heavily on it.)
- Incrementally implementing a simple idea: Best Insert.
- This consists of multiple submodules, it is due to technical or development reasons.
- Lot of time spent with debugging due to the many edge cases.
- It already includes some straightforward optimizations, but there are much room to improve.

# Results

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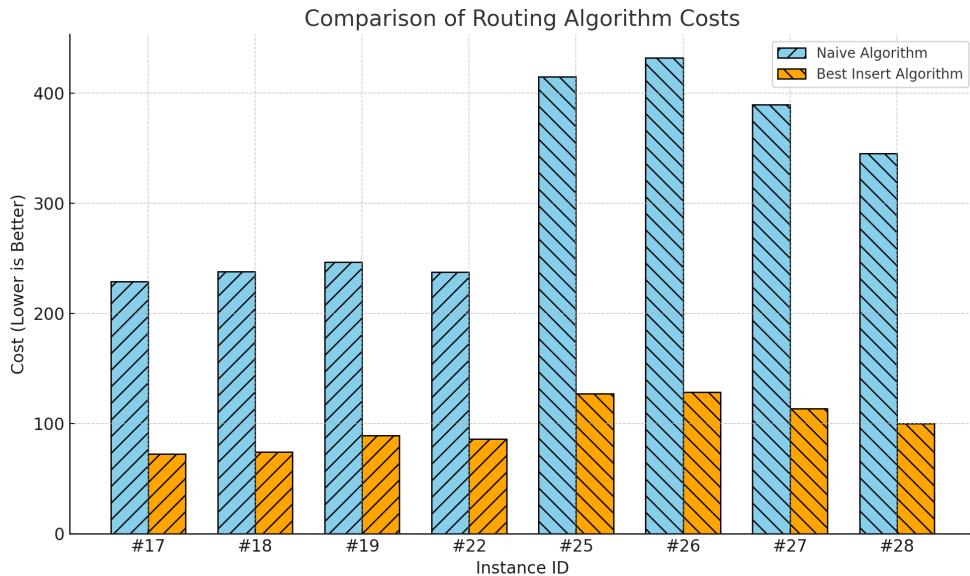


Figure: Comparison of the results between the naive algorithm and the best insert algorithm on smaller instances

# Conclusions

- Simple algorithm outperformed naive approach with **68.79% average cost reduction**.
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For instances that require waiting, further analysis and simulations are needed to explore potential improvements.

# Further research

- Implement local search methods to optimize schedules, especially early on with fewer orders.
- Explore modified objective functions (e.g., adding weighted terms) to improve flexibility and regularize solutions.
- Define intuitive policies, such as limiting vehicles per location, that are dynamically adjusted based on requests.
- Developing these ideas further as part of my master's thesis.

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Thank you for your attention!