Graph Canonization Algorithms

Nagy Szabolcs

2024/25 1st semester, Math Project 3

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ 三三 - のへぐ

Canonization overview

Isomorphism-invariant graph labeling.





う P C 三

Applications:

- **Efficient generation of non-isomorphic graphs** (n < 20),
- deciding isomorphism between graphs,
- various uses in biology in and chemistry,

etc. . .

Goal: examine and attempt to improve upon known canonization algorithms.

▲□▶ ▲□▶ ▲□▶ ▲□▶ ■ ●の00

A search-tree of ordered partitions on V(G)

- Root contains unit partition: $\Pi_0 = V(G)$.
- Partition of a child is always finer than that of the parent: $\Pi_t \subseteq \Pi_{parent(t)}$.
- Leaves are trivial partitions \iff permutations: $\Pi_{I} = \{v_{\pi_{1}}\}, \{v_{\pi_{2}}\}, \dots \{v_{\pi_{n}}\}$ for some $\pi \in S_{n}$.

Graph degree information and automorphism data are used to decrease the size of the search-tree.

A partition Π is equitable if: $\forall X_1, X_2 \in \Pi, \forall v_1, v_2 \in X_1 : d(v_1, X_2) = d(v_2, X_2)$

When setting Π_t , if it is not equitable, *refine* it so it is.

・ロト ・ 目 ・ ・ ヨト ・ ヨ ・ うへつ

Possible to do in isomorphism-invariant manner.

Equitable partition example



$$V_1 = \{A, B, C, D\}, \\ V_2 = \{E, F\}, \\ V_3 = \{G, H\}$$

◆□ > ◆□ > ◆豆 > ◆豆 > ̄豆 = のへで

Degrees: $\begin{bmatrix} 0 & 1 & 1 \\ 2 & 0 & 1 \\ 2 & 1 & 1 \end{bmatrix}$

Refinement in action



◆□▶ ◆□▶ ◆三▶ ◆三▶ ●□ ● ●

k-equitability

$$\blacksquare d_k(v,X) = |\{w \in X : \exists v \to w \text{ walk } Q : |Q| = k\}|$$

• A partition Π is k-equitable if for any $\ell = 1 \dots k$: $\forall X_1, X_2 \in \Pi, \forall v_1, v_2 \in X_1 : d_\ell(v_1, X_2) = d_\ell(v_2, X_2)$

▲□▶▲□▶▲≡▶▲≡▶ ≡ めぬる

■ Notice: 1-equitability \iff equitability

Equitable, but not 2-equitable partition



$$V_1 = \{A, B, C, D\}, \\ V_2 = \{E, F\}, \\ V_3 = \{G, H\}$$

$$d_2(A, V_3) = d_2(B, V_3) = 1$$

 $d_2(C, V_3) = d_2(D, V_3) = 2$

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへぐ

Construct Markov-chain from graph:

$$I = V(G),$$

step into any neighbor with equal probability.

Idea: calculate a stationary distribution for the above Markov-chain, get ordered partition by sorting elements based on distribution value.

This involves finding the solution of the lin. eq. sys.

 $\begin{bmatrix} P^{T} - I \\ 1 \end{bmatrix} \mu = \begin{bmatrix} \mathbf{0} \\ 1 \end{bmatrix}, \text{ where } P \text{ is the transition matrix.}$



Transition matrix:

0	0.33	0.33	0.33	0]
0.5	0	0	0.5	0
1	0	0	0	0
0.33	0.33	0	0	0.33
0	0	0	1	0

Unique stationary distribution: $\mu = \tfrac{1}{10}\{3,2,1,3,1\}$

Resulting ordered partition: $\Pi=\{3,5\},\{2\},\{1,4\}$

Problem: stationary distribution might not be unique. Enough to ensure uniqueness: irreducibility, aperiodicity. Solution: additional "connecting" vertex *z*:

$$V(G') = V(G) + z$$

•
$$zv \in E(G')$$
 for all $v \in V(G')$

Constructing the same Markov-chain for G' always yields a unique stationary distribution.

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

More information can be gleamed from the distribution by putting more thought into the construction:

 adjust probabilities based on the current partition, relations between neighboring cells (allows repeated use of distributions),

possible steps based on k-long paths between vertices (helps in "ensuring" k-equitability). Advantages of using k-equitability and stationary distributions:

- Can find finer partitions than regular equitability checking
- Decreases the number of search-tree nodes for certain graphs, reducing runtime.

Drawbacks:

- More time spent examining each node
- Search-tree not guaranteed to be reduced, many graphs are unaffected

Does not notably improve upon automorphism pruning

Canonizing all labeled graphs of size 6

37268 graphs total

- 2E 2-equitability refinement
- SD Stationary distribution
- AP automorphism pruning

Method	Method Nodes Run		# smaller trees	# bigger trees
Ø	139 995	770 ms	-	-
2E	139 935	740 ms	60	0
SD	139 770	754 ms	195	0
AP	127 915	684 ms	-	-
AP + 2E	127 915	730 ms	0	0
AP + SD	128 215	720 ms	0	90

What we did:

- maintain/optimize our canonization implementation, and the corresponding graph generator (C++)
- implement cell refinement based on 2-equitability,
- implement cell refinement based on stationary distributions.

Plans for the future:

- generalize cell refinement for k-equitability,
- further optimize Markov-chain construction,
- look into further potential improvements to canonization.

In the last 3 semesters, we have:

- thoroughly studied the theory of canonical labelings,
- comprehended the main algorithm and its implementation,
- produced our own object-oriented canonization program,
- used it to create our own customizable isomorphism-free graph generator,
- tested some potential improvements to the algorithm,
- found solid ground as to where to improve in the future.